

Model Driven Development

The process of generating software based on models through some automated form of code generation

Our Group's Work

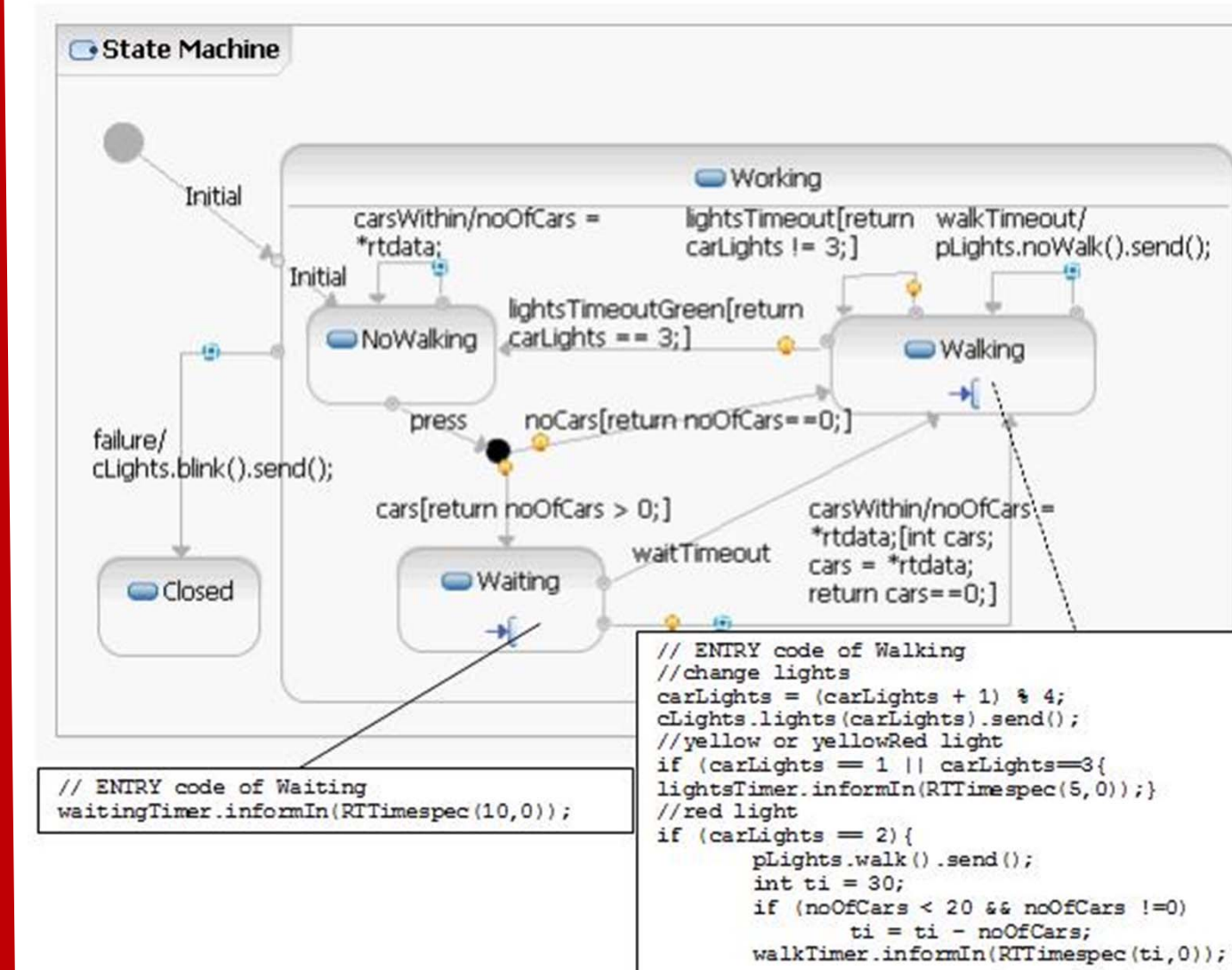
kiltera

Language for modeling and simulating concurrent, interacting real-time processes; with distribution and mobility

Why Modeling Is Important

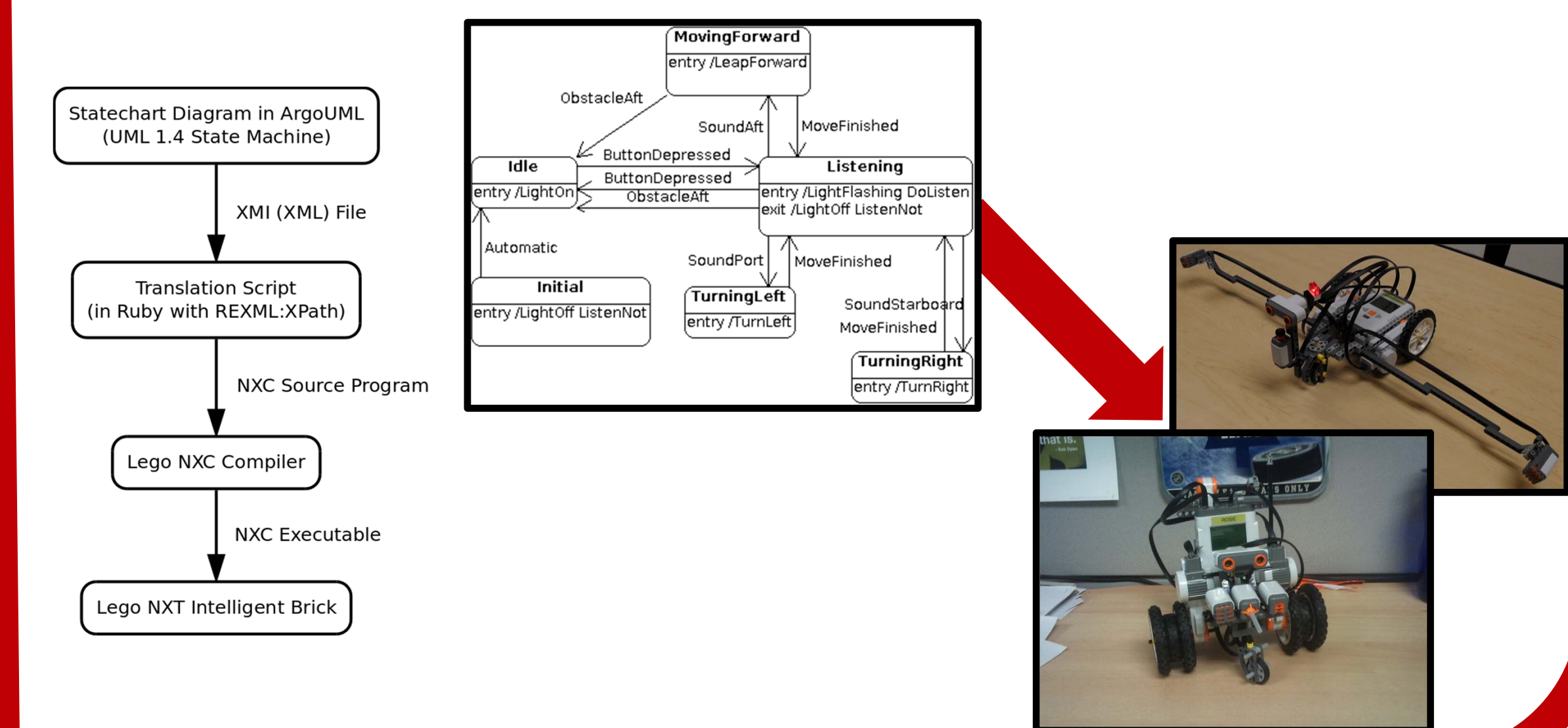
- Software has recently become a very important part of safety critical systems
- Things like missile controls, breaking systems, and others are all controlled by software
- A need exists to ensure the software is correct, but verifiable and provable languages, such as Ada, are too complex
- The solution: MDD through the use of modeling languages/toolkits such as RSA-RTE and kiltera

Rational Software Architect



- The real time edition of RSA (RSA-RTE) is used to create models of systems
- The symbolic execution can then be analyzed to determine correctness

Application: Mindstorms

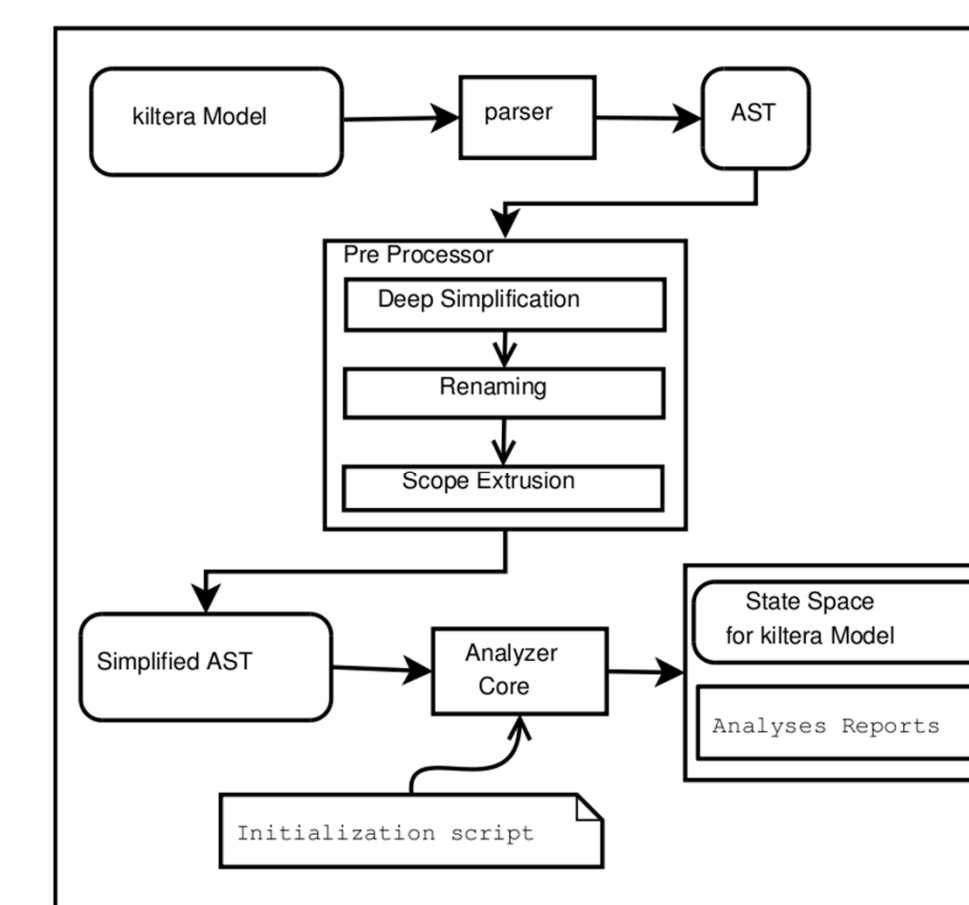


kiltera

| Name of the Class | π_{kl} Constructs |
|-------------------|--|
| PDone | nil |
| PTrig | x!E |
| PListener | when(x1?F1@y1 -> P1 x2?F2@y2 -> P2 ... xn?Fn@yn -> Pn } |
| PNew | new x in P |
| PDelay | wait E -> P |
| PPar | P1 P2 ... Pn |
| PSeq | P1 ; P2 ; ... ; Pn |
| PDef | def {D1, D2, ... , Dn} in P |
| PAssert | assert E -> P |
| PMatch | match E with { F1 -> P1 F2 -> P2 ... Fn -> Pn } |
| PCond | if E then P1 else P2 |

- Sequential and Parallel Processes
- Channels as first-class objects
- Channels can be passed between processes and PCs
- Channels can be triggered
- Channels can be listened to in parallel

Current Analyzer



- Uses the Visitor Pattern
- Parses kiltera code and builds the Abstract Syntax Tree
- Builds the State Space Tree
- Produces analyses and reports such as Deadlocks and Stable States
- Will be re-factored using the Language Definition Framework

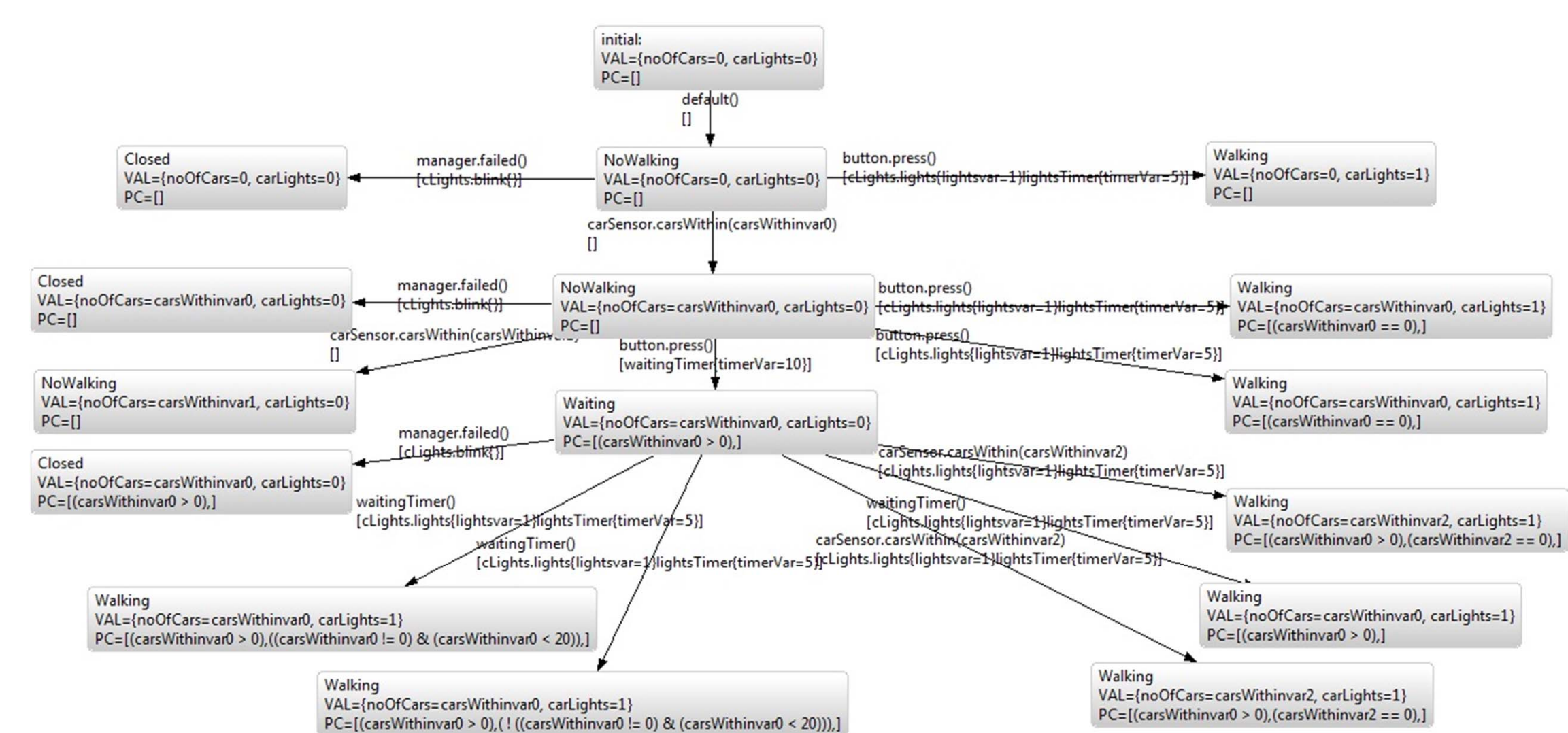
Refactoring with Language Definition

$$\begin{aligned}
 fn(\tau) &\stackrel{def}{=} \emptyset & fn(0) &\stackrel{def}{=} \emptyset \\
 fn(x) &\stackrel{def}{=} \{x\} & fn(\nu x.P) &\stackrel{def}{=} fn(P) \setminus \{x\} \\
 fn(\bar{x}) &\stackrel{def}{=} \{x\} & fn(\alpha.P) &\stackrel{def}{=} fn(\alpha) \cup fn(P) \\
 & & fn(P_1 + P_2) &\stackrel{def}{=} fn(P_1) \cup fn(P_2) \\
 & & fn(P_1 || P_2) &\stackrel{def}{=} fn(P_1) \cup fn(P_2)
 \end{aligned}$$

CCS is a Process Calculus
 $CCS \subset \pi\text{-calculus} \subset \pi_{kl} \subset \text{kiltera}$

Nicolas' Work

Symbolic Execution



Incremental Test Case Generation

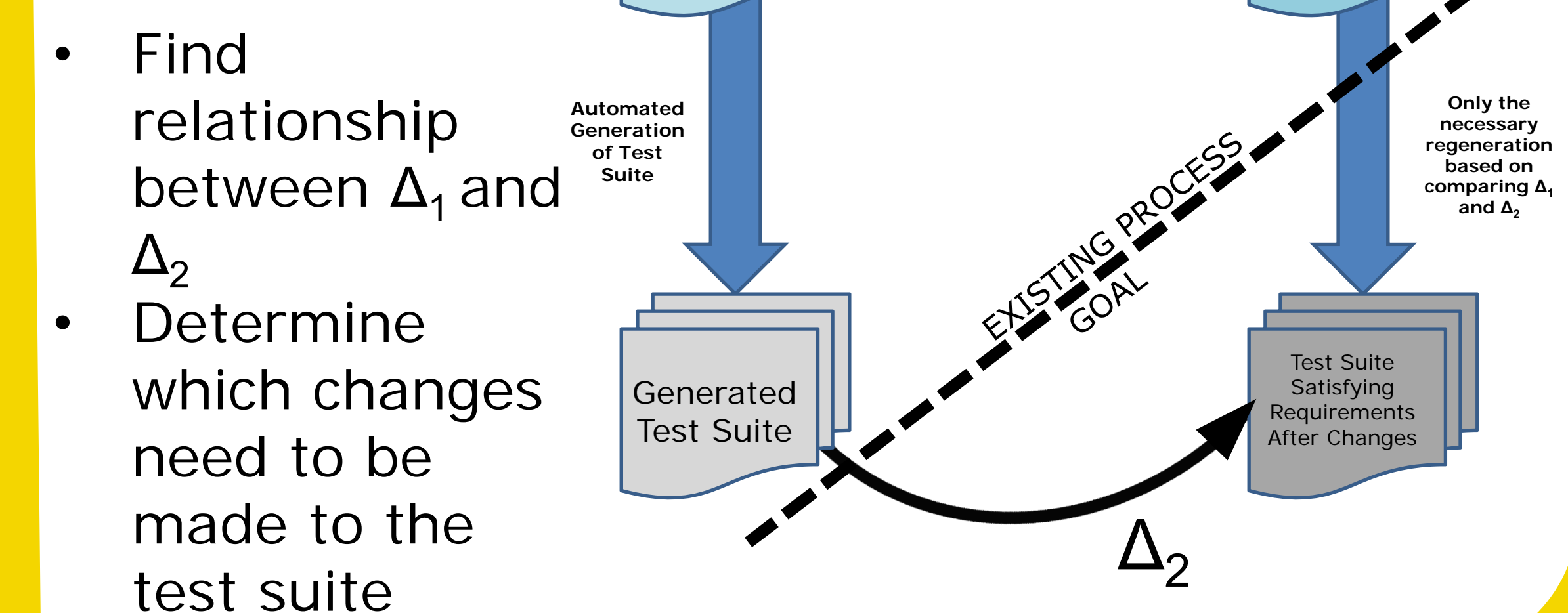
Process of automatic generation of test suites for UML-RT Models, and the effects of model changes on the generated test cases

Eric's Work

Reasons for Incremental Test Case Generation

- Automated generation of a full test suite can become a costly operation
- By determining and making only the necessary changes, this cost can be reduced
- Incrementally generating the test cases as modifications are made will ensure correctness of the test suite

The Process



- Find relationship between Δ_1 and Δ_2
- Determine which changes need to be made to the test suite

Resources

❖ Engin Uzuncaova, Sarfraz Khurshid, Don S. Batory: Incremental Test Generation for Software Product Lines. IEEE Trans. Software Eng. 36(3): 309-322 (2010)

❖ J. Dingel, E. Paen, E. Posse, R.R. Rahman, and K. Zurowska. Denition and implementation of a semantic mapping for UML-RT using a timed pi-calculus. In Proceedings of the Second International Workshop on Behaviour Modelling: Foundation and Applications, pages 1-8. ACM 2010

❖ R. Rahman. Design and Implementation of an Analyzer for a Timed π -calculus. 2010
 ❖ E. Posse. Symbolic simulation of π_{kl} . 2010